
ASM Semantics of SDL: *Concepts, Methods, Tools*

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Abstract State Machines

- Basic ASM Concepts
- Multi-Agent Real-Time ASM

Abstract SDL Machine

- Overall Organization
- Encoding of SDL Objects
- Operations on Signals and Timers

ASM Tool Support

Conclusions

ITU-T Recommendation Z.100 (\approx 200 pp. without Annexes)

Formal model of SDL (Meta-IV, CSP):

- *Annex F.2*: Static properties of SDL (437 pp.)
- *Annex F.3*: Dynamic properties of SDL (183 pp.)

Other Approaches

- φ SDL: *process algebra semantics* [Bergstra and Middleburg 96]
- Base SDL: *Object-Z* [Lau and Prinz 95]
- SDL Time Nets: *extended Petri Nets* [Fischer and Dimitrov 95]
- FOCUS: *stream processing functions* [Broy 91], [Holz and Stølen 94]
- ...

Disclaimer

Z.100, Annex F.1, p. 1:

“This annex constitutes a formal definition of SDL. If any properties of an SDL concept defined in this document, contradicts the properties defined in Z.100 and the concept is consistently defined in Z.100, then the definition in Z.100 takes precedence and this formal definition requires correction.”

Our Approach: separate *specification* from *verification*

- **operational semantics**

- ▷ embed a formal documentation into Z.100
- ▷ reflect the *common* understanding of SDL
- ▷ support alternative levels of abstraction

- **approved meta-modelling concept**

previous work:

ASM semantic of VHDL [Börger,Glässer,Müller 95]

Abstract State Machines

BASIC ASM CONCEPTS

States

Structures with domains and functions

$$f(t_1, \dots, t_n) := t_0, \quad n \geq 0 \quad \xrightarrow{\text{Update}} \quad \left\langle \underbrace{f^S(t_1^S, \dots, t_n^S)}_{\text{location}}, \underbrace{t_0^S}_{\text{value}} \right\rangle$$

Update Sets

$$\Delta_S(P) \quad \text{program } P$$

Computations

$$S_0 \xrightarrow{\Delta_{S_0}(P)} S_1 \xrightarrow{\Delta_{S_1}(P)} S_2 \xrightarrow{\Delta_{S_2}(P)} \dots \quad \text{“pure runs”}$$

General Abstraction Principles

Clear & Concise Specifications: “*ground models*”

- **abstract operational views**

- ▷ hierarchical descriptions
- ▷ incremental refinements
- ▷ *behaviour* separated from *context*

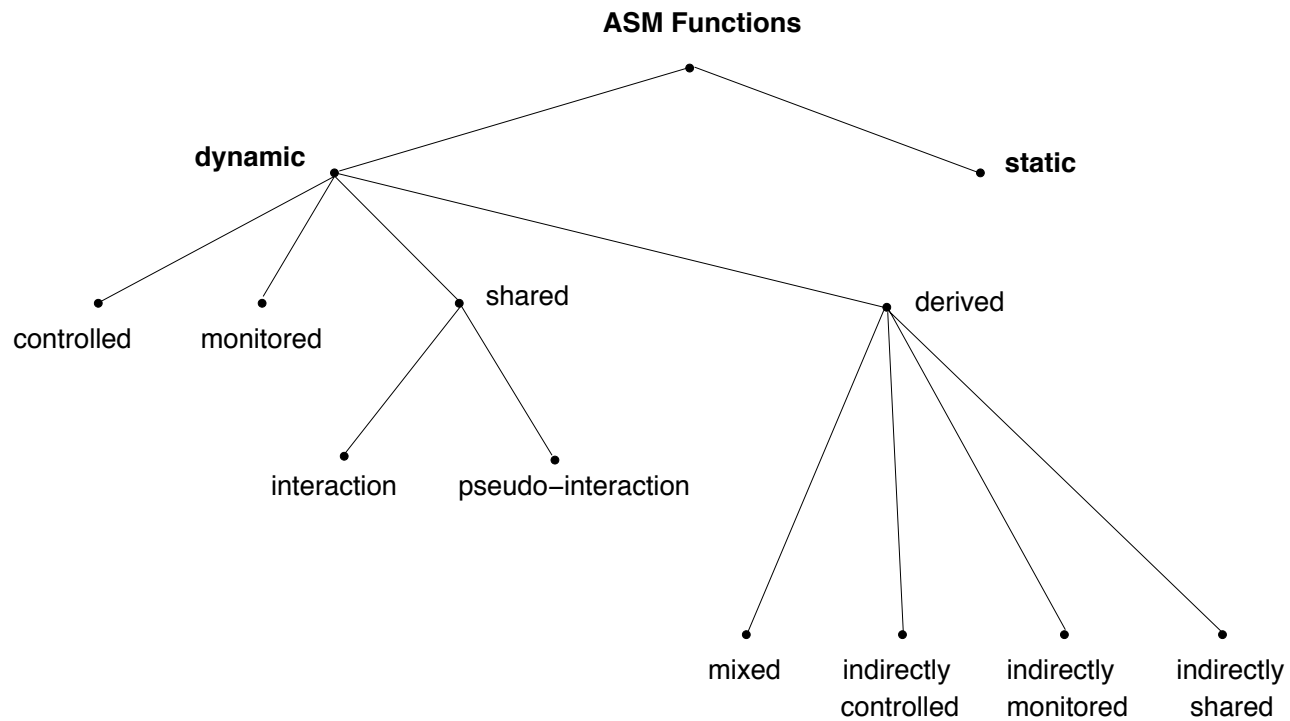
- **information hiding & interface mechanisms**

classification of ASM functions:

- static
- dynamic (*controlled, monitored, shared ... global, private*)

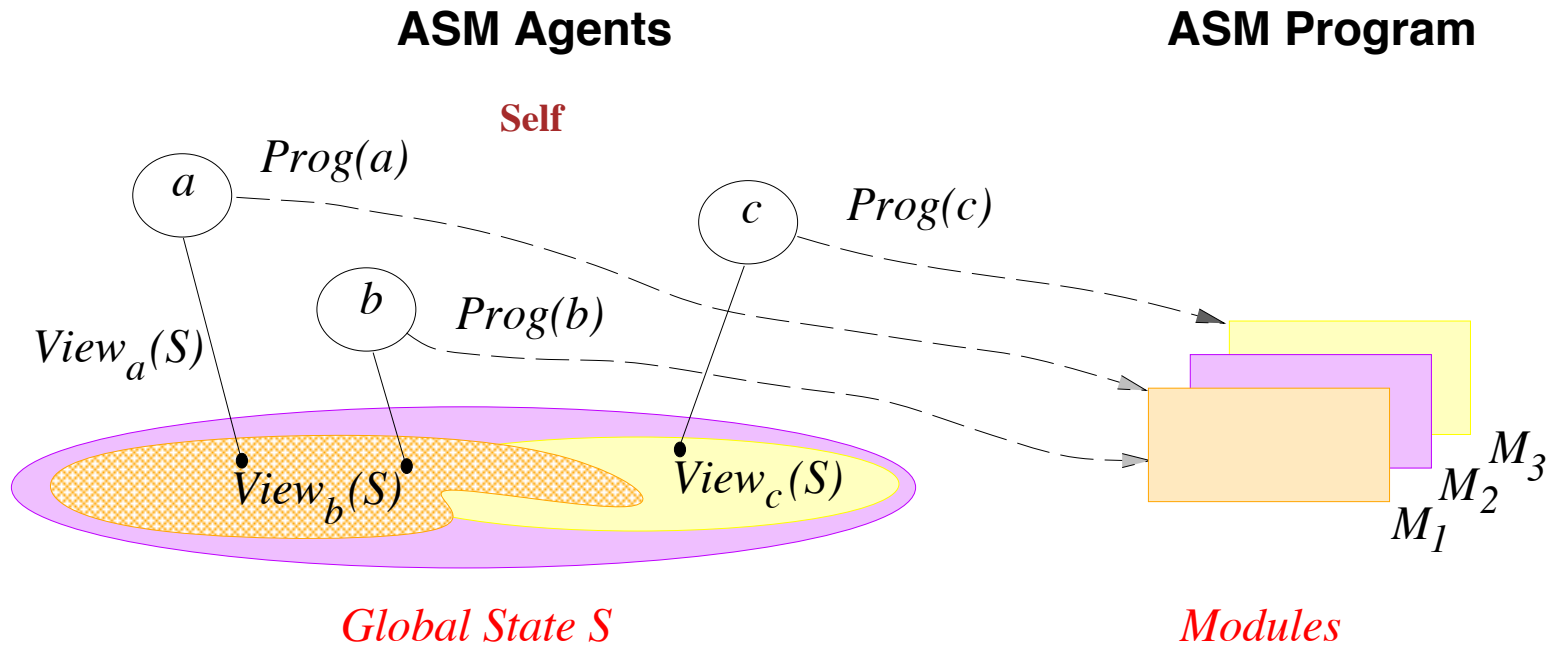
General Abstraction Principles

Classification of ASM Functions



Multi-Agent Real-Time ASM

CONCURRENCY AND NONDETERMINISM



Global System Time

external (monitored) function:

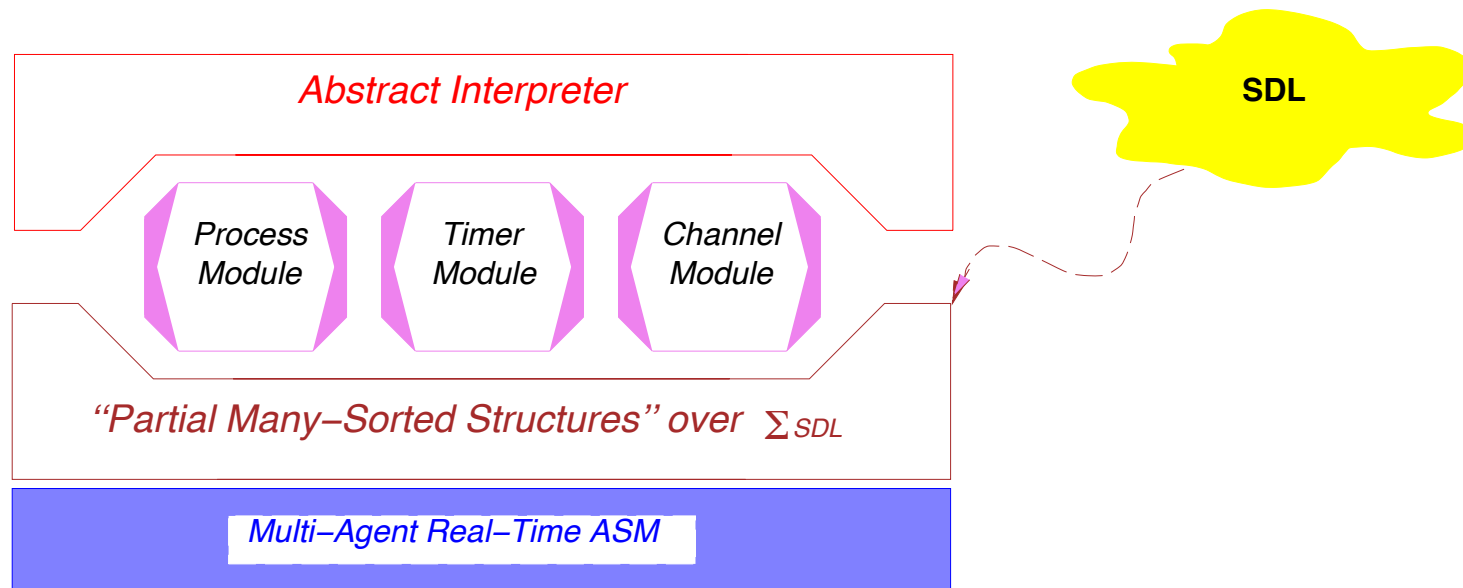
$$now : TIME, \quad TIME \subseteq \mathbb{R}$$

Agents perform *instantaneous* actions in *continuous* time:

An agent which is enabled at *time* t to fire a *rule* R actually fires R not later than $t + \epsilon$ (for some infinitely small ϵ).

Abstract SDL Machine

Interpretation Model (Basic SDL)



Encoding of SDL Objects

Static Reachability Constraints

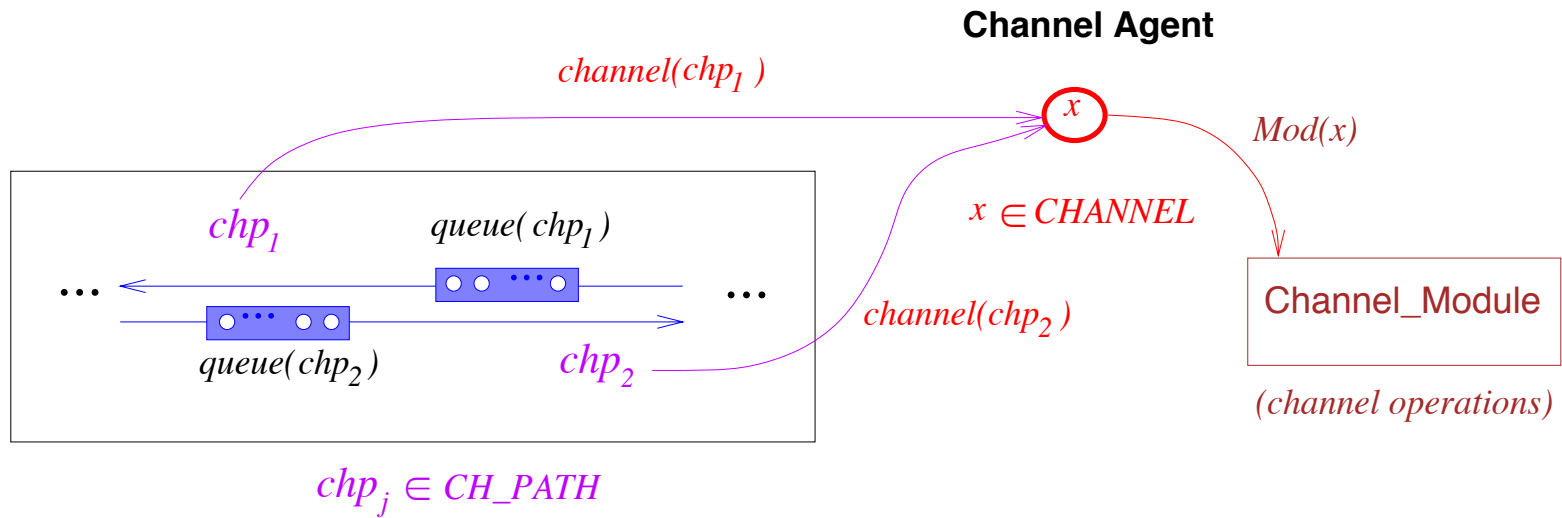
Reachability Sets

Choice of *receiver* and *path*:

$$\begin{aligned} \textit{choose_reachability} : & \textit{PROCESS} \times \textit{SIGNAL} \times \mathbf{to-Arg} \times \mathbf{via-Arg} \\ & \rightarrow \textit{PATH} \times \textit{RECEIVER} \end{aligned}$$

Encoding of SDL Objects

Bidirectional Delaying Channel



Operations on Signals

Behaviour of Channels

DELIVERSIGNALS

\equiv do forall chp : $CH_PATH(chp)$ and $channel(chp) = Self$
if $ReadyToDeliver(chp)$ then
 $queue(chp) := tail(queue(chp))$
 let $si = head(queue(chp))$, $r = receivername(si)$ in
 if $r = env$ then
 DELIVERTOENV(si)
 else
 DELIVERTOPROCESS(si, r)

where

$ReadyToDeliver(chp)$

$\equiv \exists si : SIGINST(si) \wedge si = head(queue(chp)) \wedge \neg InTransit(si, chp)$

Operations on Signals

Behaviour of Channels (continued)

$\text{DELIVERTOPROCESS}(SInst, PName)$
 $\equiv \text{let } PId = \text{receiverid}(SInst) \text{ in}$
 if $PId = \text{undef}$ then
 choose $p : PID(p)$ and $\text{procname}(p) = \text{receivername}(SInst)$
 $\text{buffer}(p) := \text{buffer}(p) \frown \langle SInst \rangle$
 else
 if $PID_{Sys}(PId)$ then
 $\text{buffer}(PId) := \text{buffer}(PId) \frown \langle SInst \rangle$

Operations on Timers

Expiration Time

$$\mathit{expire} : \mathit{TIMERINST} \rightarrow \mathit{TIME}$$

Behaviour of Timers

$$\mathit{Active}(t) \equiv$$

$$\mathit{ActiveTime}(t) \vee \mathit{ActiveSignal}(t)$$

$$\mathit{ActiveTime}(t) \equiv$$

$$\mathit{TIMERINST}(t) \wedge \mathit{expire}(t) \neq \mathit{undef}$$

$$\mathit{ActiveSignal}(t) \equiv$$

$$\mathit{TIMERINST}(t) \wedge \exists s \in \mathit{SIGINST} : t = \mathit{timer}(s) \wedge s \text{ in } \mathit{buffer}(\mathit{owner}(t))$$

Operations on Timers

Operations on Timers

TIMEROPERATION

```
≡ if MyAction(Self) then
  if Action = set then
    let time = fst(Arg) in
      SETEXPIRATIONTIME(time)
      DISCARDTIMERSIGNAL
  else
    if Active(Self) then
      expire(Self) := undef
      DISCARDTIMERSIGNAL
  else
    if ActiveTime(Self) ∧ now ≥ expire(Self) then
      expire(Self) := undef
      CREATETIMERSIGNAL
```

Operations On Timers

Behaviour of Timers (continued)

SETEXPIRATIONTIME(*Time*)

≡ if *Time* = undef then

$expire(Self) := now + duration(timename(Self))$

elif $Time \leq now$ then

$expire(Self) := undef$

 CREATETIMERSIGNAL

else

$expire(Self) := Time$

Conclusions

Observations

- *SDL view* and *ASM view* of distributed “*real-time*” systems *coincide*:
 - ▷ notions of *concurrency*, *reactivity* and *time* are tightly related;
 - ▷ common understanding of SDL can *directly* be formalized.
- ASM semantics of Basic SDL
 - ▷ is particularly *concise*, readable and understandable;
 - ▷ can easily be *extended* and *modified*;
 - ▷ *bridging semantics* for combining SDL with other languages.

Further References

<http://www.uni-paderborn.de/cs/asm/>